

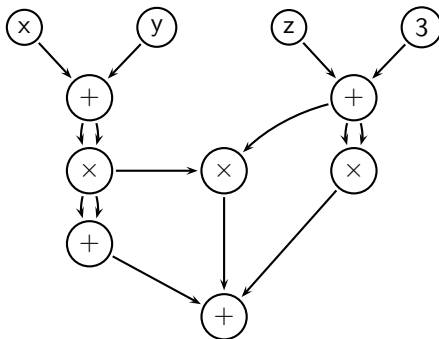
Determinant versus Permanent: salvation via generalization?

The algebraic complexity of the Fermionant and the Immanant

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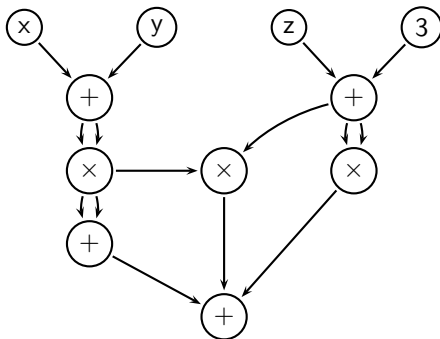
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The size of the entry is the number of variables. The *size* of an arithmetic circuit is the number of operational gates.



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A family $\mathbf{F} = (f_n)$ of polynomials is in VP if there exists a family of circuits C_n of polynomial size such that for any n

$$f_n \text{ is computed by } C_n$$

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A family $\mathbf{F} = (f_n)$ is in VNP if there is a family $\mathbf{G} = (g_n)$ in VP such that

$$f_n(\bar{x}) = \sum_{\bar{\epsilon} \in \{0,1\}^n} g_n(\bar{\epsilon}, \bar{x})$$

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Definition

Let S_n be the symmetric group on n elements and if $\pi \in S_n$, $c(\pi)$ its number of cycles. The determinant is

$$\det_n(\bar{x}) = (-1)^n \sum_{\pi \in S_n} (-1)^{c(\pi)} \prod_{i=1}^n x_{i\pi(i)}$$

Theorem (Valiant 79)

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Let S_n be the symmetric group on n elements and if $\pi \in S_n$, $c(\pi)$ its number of cycles. The permanent is

$$\text{per}_n(\bar{x}) = \sum_{\pi \in S_n} \prod_{i=1}^n x_{i\pi(i)}$$

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The family $\text{per} = (\text{per}_n)_{n \in \mathbb{N}}$ is VNP-complete

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Conjecture (Valiant hypothesis)

$$VP \neq VNP$$

Theorem (Bürgisser 2000)

Under Generalized Riemann Hypothesis,

$$VP = VNP \Rightarrow P/poly = NP/poly$$

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Three main approaches of Valiant Hypothesis:

- Geometric Complexity Theory (GCT)
- Lower bounds
- The study of complexity classes (Characterization, complete polynomials)

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Definition (informal)

A generalization of the determinant and the permanent is a series of family $\mathbf{F}^k = (f_n)^k$ indexed by some k such that

- For certain k , \mathbf{F}^k are in VP
- For others k , \mathbf{F}^k are VNP-complete.

Definition

If π is a permutation, $c(\pi)$ is its number of cycles.

$$\text{Ferm}_n^k A = (-1)^n \sum_{\pi \in S_n} (-k)^{c(\pi)} \prod_{i=1}^n A_{i, \pi(i)}$$

Let \mathbf{Ferm}^k the family of (Ferm_n^k)

- If $k = 1$, then $\text{Ferm}^1(\bar{x}) = \det(\bar{x})$
- If $k = -1$, then $\text{Ferm}^{-1}(\bar{x}) = \text{per}(\bar{x})$

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Theorem (1)

- $Ferm^0 = 0$.
- $Ferm^1$ is in VP
- for $k \in \mathbb{Q}$ different from 0, 1 $Ferm^k$ is VNP-complete.

Corollary

$Ferm^k$ is $\#P$ -complete for any rational $k \notin \{0, 1\}$.

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Lemma (Idea of the demonstration)

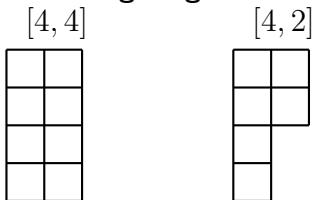
Let $X = (x_{i,j})_{i,j \in [n]}$ be the matrix of variables. For any $l \in \mathbb{N}$ there exists a transformation P_l of X such that

$$\text{Ferm}^k(P_l(X)) = \sum_{\pi \in S_n} (-k)^{c(\pi) \times l} \prod_{i=1}^n x_{i, \pi(i)}$$

Definition

A young diagram is a collection of boxes in left adjusted row with decreasing row length.

Young diagrams



Definition

Let χ_Y be an irreducible character of S_n . Then

$$\text{im}_{\chi}(\bar{x}) = \sum_{\pi \in S_n} \chi_Y(\pi) \prod_{i=1}^n x_{i, \pi(i)}$$

$$Y = \boxed{} \boxed{} \boxed{} \boxed{} \boxed{}$$

Example

If Y is a single row, then $\text{im}_Y(\bar{x}) = \text{per}(\bar{x})$.

$$Z = \begin{array}{c} \boxed{} \\ \boxed{} \\ \boxed{} \\ \boxed{} \\ \boxed{} \end{array}$$

Example

If Z is a single column, then $\text{im}_Z(\bar{x}) = \text{det}(\bar{x})$.

Theorem (Bürgisser 2000)

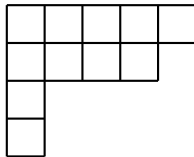
If (Y_n) is a family of Young diagrams with only a constant number of boxes at the right of the first column, then

(im_{Y_n}) is in VP

Theorem (Brylinski 2003)

Let Y_n be Young diagrams such that the maximal difference between the size of two consecutive rows is $\Omega(n)$, then

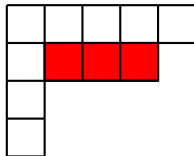
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$$[n, n] = \begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline \end{array}$$

Theorem (2)

Let $[n, n]$ be the Young diagram with two columns, each with n boxes. Then

$(im_{[n,n]})$ is VNP-complete

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Theorem (3)

If (Y_n) has a polynomial number of boxes at the right of the first column and a constant number of columns, then

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Proposition

For any integers k, n , if we write Λ_k^n the set of every Young diagrams with n boxes and at most k columns, then there exists some rational constants d_Y^k such that for any matrix A ,

$$\text{Ferm}_n^k(A) = \sum_{Y \in \Lambda_k^n} d_Y^k \text{im}_Y(A)$$

Theorem (Conclusion)

Let (Y_n) be a family of Young diagrams with a constant number of columns such that $|Y_n| = \Omega(n)$. Then

- If the number of boxes at the right of the first column is constant c , then (im_{Y_n}) is in VP.
- If the number of boxes at the right of the first column is logarithmic, then (im_{Y_n}) is not VNP-complete
- If the number of boxes at the right of the first column is polynomial, (im_{Y_n}) is VNP-complete.

Perspectives

- Studying the class of polynomials computed by sub-exponential circuits.
- Finding a VP-complete family !

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Thank you !